

# Discrete Abstraction and Supervisory Control of Switching Systems\*

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**Abstract** – *In this paper we propose a method to create discrete abstraction of state space behavior for continuous-time systems based on gradient analysis of the system dynamics. Then we describe how to use such a discrete model to design a supervisory controller for a given safety specification for the system. Finally we provide an entropy measure of nondeterminism, which can be used to evaluate the quality of the result discrete model as the degree of nondeterminism in that model. The discrete abstraction and supervisory control approach is demonstrated on a two tank system with switching control.*

**Keywords:** Discrete abstraction, supervisory control, hybrid systems.

## 1 Introduction

In this paper we consider the problem of discrete abstraction of system behavior and supervisory control for a special class of dynamic systems, referred to as *switching systems*. Switching systems are characterized by a finite set of inputs, and a finite set of operating modes. In general, hybrid systems can be described by a transition structure whose state space consists of two domains associated with the discrete-event and continuous-time dynamics. Transitions in hybrid systems are either time-based - changing the continuous-time variables according to a given set of differential/difference equations - or event-driven - changing the discrete variables and possibly resetting the continuous-time variables to a new value from which evolution is governed by another set of equations. Switching systems represents a special class of hybrid systems.

Considerable research work has been dedicated recently to the study of hybrid systems. See for example

\*0-7803-7952-7/03/\$17.00 © 2003 IEEE.

Funded, in part, by the DARPA IXO Software-Enabled Control Program under contract F33615-99-C-3611.

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[1, 2] and the references therein. In developing supervisory control of hybrid systems, discrete abstractions are usually employed to approximate the continuous dynamics of the system. As a result, the hybrid system is modelled as a discrete event system with a finite state space. Abstractions are essentially in dealing with complex systems. Consider the example of a large chemical process made up of a number of tanks, pumps, and pipes. The corrosive nature of the fluid in the tanks may make it difficult to get accurate measures of height, therefore, discrete level sensors are employed to determine if the height of fluid in a tank is above or below a pre-determined value.

From a modeling viewpoint, typical abstraction schemes first partition the state space of the system into finite regions with well-defined boundaries. States in the abstracted system corresponds to regions in the state space of the hybrid system, and events capture the crossing of the boundaries between these regions. Building approximate discrete finite state models of hybrid systems has been an active area of research for more than a decade. Several conservative abstraction techniques have been proposed in literature, e.g. in [3] methods for hybrid systems with discrete valued input and outputs; in [6, 7] for linear systems are presented; in [8, 11, 4] for nonlinear continuous-time systems are given; and in [9, 10] methods for nonlinear discretizable systems are proposed.

The proposed technique is aimed at reducing a switching system into a finite state structure that preserves the important dynamics of the original system. It addresses the task of discrete abstraction by performing gradient analysis at the boundaries of different cells in a chosen partition on the state space of the continuous-time system. Since computation is restricted to boundary regime, this method is more computationally efficient than trajectory-based abstraction techniques in the literature, e.g., [6, 9, 10]. Abstraction based on gradient analysis has been proposed in [8, 4]. However, these approaches target a general setting with no specific assumptions about the target system, and, therefore, did not yield a concrete computation procedure for gradi-

ent analysis. The abstraction problem is addressed and solved in detail in this paper. A supervisory control scheme for the abstracted system is then proposed. Finally, we also propose a method to measure nondeterminism of a discrete model, which offers us a way to evaluate the quality of the given partition.

This paper is organized as follows. In section 2, we describe how to apply gradient analysis in discrete abstraction for a general class of switching systems. In section 3 we describe the application of supervisory control technique on the discrete model. In section 4 we discuss how to measure nondeterminism of a discrete-event mode. Conclusions and further research directions are presented in section 5.

## 2 Discrete abstraction

Given an  $n$ -dimensional continuous-time dynamic system, which consists of  $p$  discrete modes, the dynamic model in a mode  $M_i$  takes on the following form,

$$\dot{x} = f_{M_i}(x, u) \quad \text{with } x \in X_i, \quad (1)$$

where  $x = [x_1, \dots, x_n]$ ,  $f_{M_i} = [f_1, \dots, f_n]$  and  $u = [u_1, \dots, u_m]$  are vectors. As a simple example, we use the two-tank system shown in Figure 1. It consists of two tanks ( $T_1$  and  $T_2$ ) and two valves ( $V_1, V_2$ ). Valve operations are binary, i.e., they can be fully open (=1) or fully close (=0). When  $V_1$  is open, there is a steady water flow into the system. Water level in each tank is monitored by a level-crossing sensor, which only reports whether water level is above or below some predefined values, e.g. above 0.3m or below 0.3m. The state-space

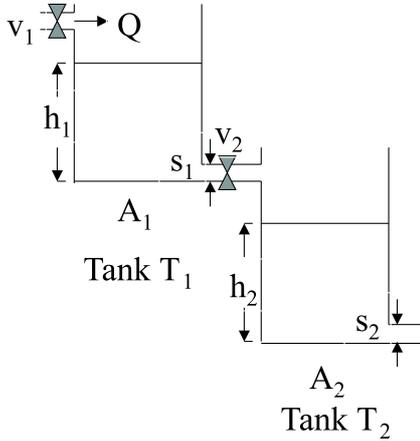


Figure 1: Two-Tank System

equations for the system, shown below cover all operating modes:

$$\dot{x} = \begin{bmatrix} \dot{h}_1 \\ \dot{h}_2 \end{bmatrix} = \begin{bmatrix} \frac{u_1 Q - u_2 s_1 \sqrt{2gh_1}}{A_1} \\ \frac{u_2 s_1 \sqrt{2gh_1} - s_2 \sqrt{2gh_2}}{A_2} \end{bmatrix} = f(x, u) \quad (2)$$

where  $Q = 5m^3/s$  is the steady water flow rate,  $h_i \in [0, 9)(m)$ ,  $A_i = 2m^2$ ,  $s_i = 0.3m^2$  is the intersection area of Tank  $i$ ,  $u_i \in \{0 : \text{off}, 1 : \text{on}\}$  is the control action of  $V_i$  with  $i \in [1, 2]$ , and  $g = 9.8Nm/s^2$  is the gravity constant

Given a set  $X$ , we use the notation  $\partial X$  to mean the boundary of  $X$  (i.e., the set of limit points of  $X$ ). An  $n$ -dimensional polyhedron  $X \subseteq \mathcal{R}^n$  is *regular* if there is a set of lower and upper bounds  $\{a_{i,l}, a_{i,h} \in \mathcal{R} | a_{i,l} < a_{i,h} \wedge 1 \leq i \leq n\}$  such that

$$X \cup \partial X = \{(x_1, \dots, x_n) \in \mathcal{R}^n | a_{i,l} \leq x_i \leq a_{i,h}, i \in [1, n]\}.$$

We make the following assumptions about the continuous-time dynamic system:

**A<sub>1</sub>:**  $u$  is from a finite set  $U$ .

**A<sub>2</sub>:** For each  $i$ ,  $X_i$  is a regular polyhedron, and so is  $X = \bigcup_{i=1}^m X_i$ .

**A<sub>3</sub>:**  $(\forall i, j) X_i \cap X_j = \emptyset$ .

**A<sub>4</sub>:** For each control action  $u \in U$ ,  $\dot{x}$  is continuous in  $X \cup \partial X$ .

With those assumptions our objective is to build a discrete-event model

$$\mathbf{G} = (Z, U, \xi),$$

where  $Z$  is the *state set*,  $U$  is the *transition event set*, and  $\xi : Z \times U \rightarrow Z$  is the (*partial*) *transition function*.  $\mathbf{G}$  is supposed to be an *abstraction* (which will be defined latter) of the continuous-time system.

Assume that for each  $i$ ,  $X_i$  is partitioned into a set  $Z_i$  of regular  $n$ -dimensional polyhedrons. Then  $Z = \bigcup_{i=1}^m Z_i$ . In the two-tank system,  $X = \{(h_1, h_2) | h_1 \in [0, 9) \wedge h_2 \in [0, 9)\}$ , hence it is a regular polyhedron. We divide the height of each tank into three equal intervals  $[0, 3)$ ,  $[3, 6)$ ,  $[6, 9)$  and let  $W$  be the set of these intervals. Then  $Z = W \times W$  and each element  $z \in Z$  is a regular polyhedron. Figure 2 depicts the partition on  $X$ .

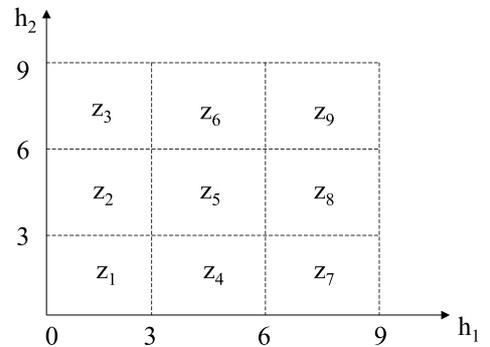


Figure 2: Partition on  $X$  in Two-Tank System

It is easy to check that the two-tank system model satisfies all above assumptions. Let  $T_{u,x} : [0, \infty) \rightarrow X$  be a trajectory of the continuous-time system under control action  $u$  with  $T_{u,x}(0) = x$ . Let  $T = \{T_{u,x} | u \in U \wedge x \in X\}$  be the set of all possible trajectories.

**Definition 2.1** A discrete-event model  $\mathbf{G}$  is an *abstraction* of the continuous-time model  $\dot{x} = f(x, u)$ , if for each pair of states  $z_1, z_2 \in Z$  and each control action  $u \in U$ ,  $z_2 \in \xi(z_1, u)$  if and only if there exists a trajectory  $T_{u,x} \in T$  such that

1.  $T_{u,x}(0) \in z_1$  and  $T_{u,x}(t) \in z_2$  for some  $t < \infty$ ,
2.  $\{t' | 0 \leq t' \leq t \wedge T_{u,x}(t') \in z_1 \cup z_2\}$  is dense in  $[0, t]$ .

□

The second condition means that “almost all” parts of  $T_{u,x}$  is in  $z_1 \cup z_2$ ; if it does include points which are not in  $z_1 \cup z_2$ . Such “almost all” requirement instead of “all” has practical reasons which will be made clear later. Notice that Def. 2.1 implies that there exists a transition between  $z_i$  and  $z_j$  only if they are neighbored to each other. Otherwise by assumption of continuity  $\mathbf{A}_4$ , condition 2 in Def. 2.1 won't hold. Thus if such trajectory  $T_{u,x}$  does exist then it must cross the *boundary* of  $z_1, z_2$ , namely  $\bar{z}_1 \cap \bar{z}_2$  where the symbol  $\bar{z}_1$  represents the closure of  $z_1$  (i.e.  $z_1 \cup \partial z_1$ ). So in order to build the (partial) transition function  $\xi$  we only need to analyze those boundary crossings. Given two states:

$$\begin{aligned} z_i &= \{(x_1, \dots, x_n) \in X | (\forall k) x_k \in [b_{k,l}^i, b_{k,h}^i]\} \\ z_j &= \{(x_1, \dots, x_n) \in X | (\forall k) x_k \in [c_{k,l}^j, c_{k,h}^j]\} \end{aligned}$$

Let  $d_{k,l}^{ij} = \max\{b_{k,l}^i, c_{k,l}^j\}$  and  $d_{k,h}^{ij} = \min\{b_{k,h}^i, c_{k,h}^j\}$ . Then their boundary is defined by

$$\begin{aligned} B(z_i, z_j) &= \{(x_1, \dots, x_n) \in X | \\ &(\forall k) x_k \in [d_{k,l}^{ij}, d_{k,h}^{ij}] \wedge d_{k,l}^{ij} \leq d_{k,h}^{ij}\} \end{aligned}$$

If  $B(z_i, z_j) = \emptyset$  then states  $z_i$  and  $z_j$  are not adjacent to each other, hence no transition exists between them. Now suppose  $B(z_i, z_j) \neq \emptyset$  and  $z_i$  in mode  $M_i$ ,  $z_j$  in mode  $M_j$ . We discuss how to determine boundary crossings on  $B(z_i, z_j)$ .

By assumption  $\mathbf{A}_3$  we get that  $\dim(B(z_i, z_j)) < n$ , namely  $(\exists k) d_{k,l}^{ij} = d_{k,h}^{ij}$ . We call  $x_k = d_{k,l}^{ij} = d_{k,h}^{ij}$  a *border*. Let  $\mathcal{I}(z_i, z_j) = \{k | d_{k,l}^{ij} = d_{k,h}^{ij}\}$  and

$$\begin{aligned} \phi(z_i, z_j) &:= \{(x_1, \dots, x_n)_k | k \in \mathcal{I}(z_i, z_j) \wedge \\ x_k &= \frac{c_{k,l}^j + c_{k,h}^j - b_{k,l}^i - b_{k,h}^i}{|c_{k,l}^j + c_{k,h}^j - b_{k,l}^i - b_{k,h}^i|} \wedge (\forall q) q \neq k \Rightarrow x_q = 0\} \end{aligned}$$

By assumptions  $\mathbf{A}_2$  both  $z_i$  and  $z_j$  are regular polyhedrons, so  $b_{k,l}^i < b_{k,h}^i$  and  $c_{k,l}^j < c_{k,h}^j$ . Considering that  $\max\{b_{k,l}^i, c_{k,l}^j\} = \min\{b_{k,h}^i, c_{k,h}^j\}$ , we get either

$b_{k,h}^i > b_{k,l}^i = c_{k,h}^j > c_{k,l}^j$  or  $c_{k,h}^j > c_{k,l}^j = b_{k,h}^i > b_{k,l}^i$ . In either case,  $\phi(z_i, z_j)$  is well defined as a set of unit vectors. We use the notation “ $\cdot$ ” to represent the dot product of two vectors.

**Proposition 2.1** Given two neighboring states  $z_i, z_j \in Z$  and control action  $u \in U$ , if

$$(\exists \hat{x} \in B(z_i, z_j)) (\forall \vec{r} \in \phi(z_i, z_j)) f_{M_i}(\hat{x}, u) \cdot \vec{r} > 0 \quad (3)$$

then there exists a trajectory  $T_{u,x} \in T$  with  $T_{u,x}(0) \in z_i$  and  $T_{u,x}(t) \in z_j$  for some  $t < \infty$ , and the set  $\{t' | 0 \leq t' \leq t \wedge T_{u,x}(t') \in z_i \cup z_j\}$  is dense in  $[0, t]$ .

(Sketch) Proof: If such  $\hat{x}$  exists then by assumption  $\mathbf{A}_4$  there exists  $\tilde{x}$  in the neighborhood of  $\hat{x}$  such that (1)  $\tilde{x} \in B(z_i, z_j) \setminus \partial X$ ; (2)  $(\forall \vec{r} \in \phi(z_i, z_j)) f_{M_i}(\tilde{x}, u) \cdot \vec{r} > 0$ . Since  $\tilde{x}$  is an internal point of  $X$  and  $\hat{x}$  is continuous anywhere in  $X$ , we get that

$$(\exists \tilde{T}_{u,x} \in T) (\exists t) \tilde{T}_{u,x}(t) = \tilde{x}$$

Again, by continuity we get that

$$\begin{aligned} \lim_{\Delta t \rightarrow 0} \frac{\tilde{T}_{u,x}(t) - \tilde{T}_{u,x}(t - \Delta t)}{\Delta t} &= f_{M_i}(\tilde{x}, u) = \\ f_{M_j}(\tilde{x}, u) &= \lim_{\Delta t \rightarrow 0} \frac{\tilde{T}_{u,x}(t + \Delta t) - \tilde{T}_{u,x}(t)}{\Delta t} \quad (4) \end{aligned}$$

Therefore, there exists  $\epsilon > 0$  such that  $\Delta t < \epsilon$  implies

$$\begin{aligned} (\forall \vec{r} \in \phi(z_i, z_j)) (\tilde{T}_{u,x}(t) - \tilde{T}_{u,x}(t - \Delta t)) \cdot \vec{r} &> 0 \wedge \\ (\tilde{T}_{u,x}(t + \Delta t) - \tilde{T}_{u,x}(t)) \cdot \vec{r} &> 0 \end{aligned}$$

If we make  $\Delta t$  small enough, then by the definitions of  $\phi(z_i, z_j)$ ,  $B(z_i, z_j)$  and the fact that  $\tilde{T}_{u,x}(t) = \tilde{x} \in B(z_i, z_j)$ , we get that

$$(\forall \Delta t') \Delta t' \leq \Delta t \Rightarrow \tilde{T}_{u,x}(t - \Delta t') \in z_i \wedge \tilde{T}_{u,x}(t + \Delta t') \in z_j$$

Therefore we can define a trajectory  $T_{u,x'} \in T$  with

$$(\forall t' : 0 \leq t' \leq 2\Delta t) T_{u,x'}(t') := \tilde{T}_{u,x}(t - \Delta t + t')$$

Clearly  $T_{u,x'}(0) = x' = \tilde{T}_{u,x}(t - \Delta t) \in z_i$ ,  $T_{u,x'}(2\Delta t) = \tilde{T}_{u,x}(t + \Delta t) \in z_j$  and

$$[0, 2\Delta t] \setminus \{\Delta t\} \subseteq \{t' \in [0, 2\Delta t] | T_{u,x'}(t') \in z_i \cup z_j\} \subseteq [0, 2\Delta t]$$

Hence the proposition is true. □

The existence of such  $\hat{x}$  is a sufficient condition for the transition  $z_j \in \xi(z_i, u)$ , but not necessary in general. Its necessary condition depends on higher order (left-hand side and right-hand side) derivatives of  $x(t)$  on  $t$ . If we interpret  $x(t)$  as a displacement, then  $f_{M_i}(x, u) \cdot \vec{r}$  is the velocity on the direction  $\vec{r}$ . Even if  $f_{M_i}(x, u) \cdot \vec{r} = 0$ , namely the velocity projected onto  $\vec{r}$  (i.e., the normal direction of a border) is zero, the trajectory can still

cross the border from the current point instantaneously as long as the acceleration ( $\ddot{x}$ ) or an even higher but finite order of derivative projected onto  $\vec{r}$  is nonzero at the current point. If the high derivatives don't exist, then we need to consider left-hand side and right-hand side derivatives at the point where the velocity is zero. Considering the complexity involved in computing high order derivatives, we will not discuss it in this paper. Now the problem is to determine whether or not such  $\hat{x}$  in Condition (3) exists. This can be converted to the following optimization problem.

$$\begin{aligned} & \text{minimize} && \mathbf{J}_{z_i, z_j, u} = \sum_{\vec{r} \in \phi(z_i, z_j)} (f_{M_i}(x, u) \cdot \vec{r} - y_{\vec{r}})^2 \\ & \text{subject to} && x \in B(z_i, z_j) \quad \text{and} \\ & && y = \{y_{\vec{r}} \geq 0 \mid \vec{r} \in \phi(z_i, z_j)\} \end{aligned}$$

We say  $y > 0$  if  $(\forall \vec{r} \in \phi(z_i, z_j)) y_{\vec{r}} > 0$ .

**Proposition 2.2** Condition (3) in Proposition 2.1 holds iff  $y > 0$  and  $\mathbf{J}(z_i, z_j, u) = 0$ .

Proof: For the ‘‘only if’’ part, if

$$(\exists \hat{x} \in B(z_i, z_j)) (\forall \vec{r} \in \phi(z_i, z_j)) f_{M_i}(\hat{x}, u) \cdot \vec{r} > 0$$

then by **A<sub>4</sub>**,  $f_{M_i}(\hat{x}, u) \cdot \vec{r} < \infty$ . Let  $y_{\vec{r}} = f_{M_i}(\hat{x}, u) \cdot \vec{r} > 0$  and this makes  $\mathbf{J}(z_i, z_j, u) = 0$ .

For the ‘‘if’’ part,

$$\mathbf{J}(z_i, z_j, u) = 0 \Rightarrow (\forall \vec{r} \in \phi(z_i, z_j)) f_{M_i}(\hat{x}, u) \cdot \vec{r} - y_{\vec{r}} = 0$$

Since  $y > 0$ , we finally get that

$$(\forall \vec{r} \in \phi(z_i, z_j)) f_{M_i}(\hat{x}, u) \cdot \vec{r} > 0$$

By Prop. 2.1 we have  $z_j \in \xi(z_i, u)$ , as required.  $\square$

Solvers for the above nonlinear constraint minimization problem can be found in several softwares, e.g. Optimization Toolbox in MATLAB, MathOptimizer in Mathematica. A discrete abstraction procedure can be described as follows. Suppose the resultant discrete model is stored in a matrix *DiscreteModel*(*ControlAction*, *TargetState*, *ExitState*), where *DiscreteModel*( $u, z_j, z_i$ ) = 1 means there is a transition  $u$  from  $z_i$  to  $z_j$ , namely  $z_j \in \xi(z_i, u)$ . Initially all entries in *DiscreteModel* are set to be zero.

**Discrete abstraction procedure:** Suppose  $Z = \{z_1, \dots, z_k\}$  and  $U = \{u_1, \dots, u_r\}$ .

1. for  $i = 1$  to  $k - 1$  {
2.   for  $j = i + 1$  to  $k$  {
3.     for  $l = 1$  to  $r$  {
4.      if  $\mathbf{J}_{z_i, z_j, u_l} = 0$  and  $y_{z_i, z_j, u} > 0$ ,
5.        *DiscreteModel*( $u, z_j, z_i$ ) := 1;
6.      if  $\mathbf{J}_{z_j, z_i, u_l} = 0$  and  $y_{z_i, z_j, u} > 0$ ,
7.        *DiscreteModel*( $u, z_i, z_j$ ) := 1;}}

In the two-tank system, after we put in all values of parameters, the system dynamic model is as follows,

$$\dot{x} = \begin{bmatrix} \dot{h}_1 \\ \dot{h}_2 \end{bmatrix} = \begin{bmatrix} 2.5u_1 - 0.89u_2\sqrt{h_1} \\ 0.89u_2\sqrt{h_1} - 0.49\sqrt{h_2} \end{bmatrix} = f(x, u) \quad (5)$$

Assume  $z_i = [0, 3] \times [6, 9)$  and  $z_j = [3, 4) \times [6, 9)$ . Then  $B(z_i, z_j) = [3, 3] \times [6, 9)$ ,  $\mathcal{I}(z_i, z_j) = \{1\}$  and  $\phi(z_i, z_j) = \{(1, 0)\}$ . Suppose the control action  $u = (1, 0)$ , namely  $V_1$  is open and  $V_2$  is close. Then we get that

$$\begin{aligned} \mathbf{J}(z_i, z_j, u) &= \min_{x, y} \sum_{\vec{r} \in \phi(z_i, z_j)} (f_{M_j}(x, u) \cdot \vec{r} - y_{\vec{r}})^2 \\ &= \min(2.5 - y_{(1,0)})^2 \end{aligned}$$

subject to  $x \in B(z_i, z_j)$  and  $y_{(1,0)} > 0$ . Clearly  $\mathbf{J}(z_i, z_j, u) = 0$ . So by Propositions 2.1 and 2.2,  $z_j \in \xi(z_i, u)$ . Suppose  $z_j = [3, 6) \times [3, 6)$ . Then  $B(z_i, z_j) = [3, 3] \times [6, 6) = \{(3, 6)\}$ . This gives that  $\phi(z_i, z_j) = \{r_1 = (1, 0), r_2 = (0, -1)\}$ . With the same control action  $u = (1, 0)$  we have

$$\begin{aligned} \mathbf{J}(z_i, z_j, u) &= \min \sum_{\vec{r} \in \phi(z_i, z_j)} (f_{M_j}(x, u) \cdot \vec{r} - y_{\vec{r}})^2 \\ &= \min[(2.5 - y_{(1,0)})^2 + (1.2 - y_{(0,-1)})^2] \end{aligned}$$

Again  $\mathbf{J}(z_i, z_j, u) = 0$  which, by Propositions 2.1 and 2.2, implies that  $z_j \in \xi(z_i, u)$ . In the second case the boundary point  $x = (3, 6)$  is not in  $z_i \cup z_j$ . But there is a trajectory  $T_{u,x}$  with  $T_{u,x}(0) = (2.8, (\sqrt{6+0.0784})^2) \in z_i$  and  $T_{u,x}(0.1) = (3.05, 5.904) \in z_j$  such that  $\{t' \mid 0 \leq t' \leq 0.1 \wedge T_{u,x}(t') \in z_i \cup z_j\} = [0, t'] \setminus \{0.08\}$ , which is obviously dense in  $[0, t]$ . This trajectory passes the point  $x = (3, 6)$  at  $t = 0.08$  instantaneously. Although in general justifying the correctness of assumption **A<sub>5</sub>** for a continuous-time system may be very difficult, we can analytically show that it holds in the two-tank system. So we can apply the above procedure to do discrete abstraction. Figure 3 displays the result discrete-event model. In this picture the value in each state is the state number and the value on each edge is the control action. Control actions are defined in Table 1 and state numbers are defined in Table 2.

Table 1: Control action for the Two tanks System

Control Action	$u_1$	$u_2$
1	0 ( $V_1$ close)	0 ( $V_2$ close)
2	0 ( $V_1$ close)	1 ( $V_2$ open)
3	1 ( $V_1$ open)	0 ( $V_2$ close)
4	1 ( $V_1$ open)	1 ( $V_2$ open)

The above discrete abstraction procedure has also been applied to the well-known three-tank system. Once we have a discrete model  $\mathbf{G} = (Z, U, \xi)$ , we can impose supervisory control which is described in the next section.



that is not in  $Z_m$ . Since  $z(0) = z_0 \in Z_m$  and  $\theta(z_0) \neq \epsilon_o$ , there exists  $j < k$  such that

$$z_j \in Z_m \& \theta(z_j) \neq \epsilon_o \& (\forall r \in \mathbb{N}) 1 \leq r < k - j \Rightarrow z(j+r) \in Z_m \& \theta(z(j+r)) = \epsilon \quad (6)$$

But this is impossible because by the control map,  $u(j) = C(z(j), u(j-1)) \in V(z(j))$ , which implies that (6) cannot happen.  $\square$

In the two-tank system we partition the desirable zone  $[3, 6] \times [3, 6]$  as shown in Figure 4.

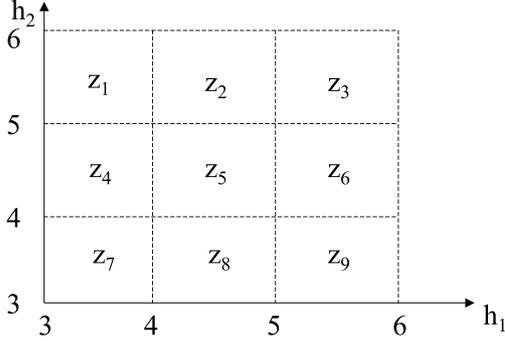


Figure 4: Desirable States for Two-Tank System

By the proposed abstraction method we obtain the following discrete model as in Figure 5, where  $a_1 = [u_1 = 0, u_2 = 0]$ ,  $a_2 = [u_1 = 1, u_2 = 0]$ ,  $a_3 = [u_1 = 0, u_2 = 1]$  and  $a_4 = [u_1 = 1, u_2 = 1]$ . Circles with dark colors represents states in  $Z \setminus Z_m$ .

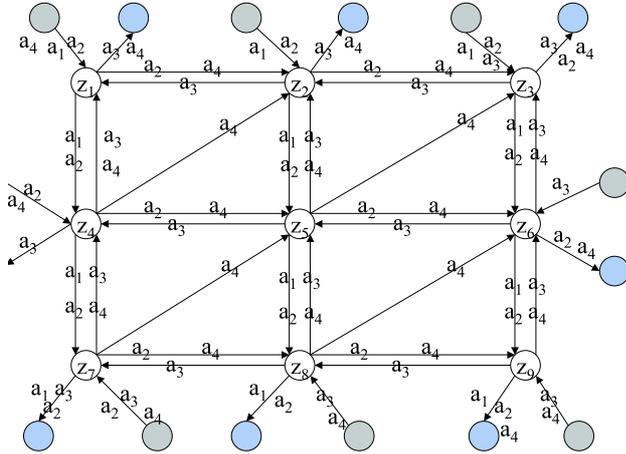


Figure 5: Discrete Model for Two-Tank System in Desirable Zone

Suppose the output event set  $\Sigma_O = Z$ , namely each state generates an observable output. If the initial state is in  $Z_m = \{z_1, \dots, z_9\}$ , then the close loop controlled behavior is shown in Figure 6.

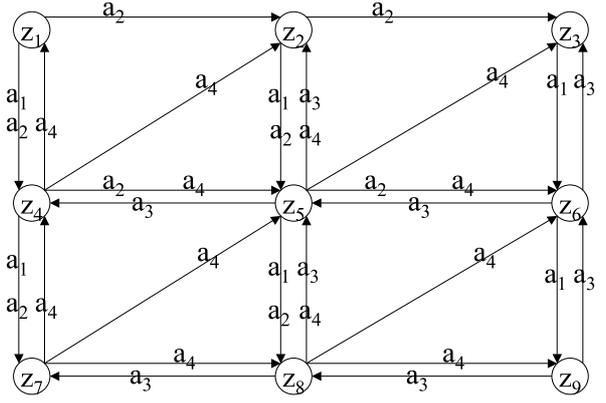


Figure 6: Close Loop Model for Two-Tank System

So far we have seen that, even though the proposed gradient-based method can only produce a rough model, it may still be able to fulfil control task as shown in the two-tank system.

## 4 Measure of nondeterminism in discrete-event models

Usually a discrete-event model generated from discrete abstraction is nondeterministic (see Figure 3), namely at some state one transition could end up in a set of states instead of one state. Nondeterminism in a discrete model is caused by the partition on the state space. In this paper we propose a measurement of nondeterminism. If the discrete model is deterministic then given the initial state  $z_0$ , a specific input sequence always generates the same sequence of states no matter how many times we run that input sequence. However, if the model is nondeterministic then with the same input sequence, we may get difference sequences of states in different runs. When the total number of test runs is very large, the frequency of each specific sequence of states will approach a fixed number, which is the probability of the occurrence of that specific sequence of states under the same input sequence.

Intuitively we can see that if some sequence of states has very high probability and the rest sequences have relatively low probability, then we may say that the model is very close to a deterministic model. In other words, the higher the difference among probabilities of different sequence of states under the same input sequence, the closer the model to a deterministic one. This intuition suggests us that we should consider each possible input sequence and compute the occurrence probability of each possible output sequence. But clearly this way is not easy to follow since we may have an infinite number of input sequences. So here we take a “weak” version of the above intuition. Under the same input sequence  $u_1 \dots u_n \in U^*$ , a sequence

of states  $z_1, \dots, z_n \in Z$  has very high probability of occurrence if for each  $i$  ( $1 \leq i \leq n$ ),  $\xi(z_{i-1}, u_i)$  has very high probability to reach  $z_i$ , namely the prior probability  $p(u_i, z_{i-1}|z_i)$  should be very high. Although the product of locally high probabilities may not lead to globally high probability, it is almost true when at each state the difference of local prior probabilities is very high. So we can say that, if at each state  $z$  and for each input event  $u$ , transition  $\xi(z, u)$  is more prone to a specific state  $z'$  than to any other state  $z'' \in \xi(z, u)$ , then the model is more close to a deterministic model. As we know, entropy is a very good measure about such a bias among transitions with the same exit state  $z$  and the same event  $u$ . So we apply it as follows.

First, we measure the nondeterminism at each state. Assume that we have a prior probability distribution  $p : U \times Z \times Z \rightarrow [0, 1]$  on each transition, namely  $p(u, z'|z)$  is defined if  $z' \in \xi(z, u)$ . Hence

$$(\forall z \in Z) \sum_{\substack{u, z', \\ z' \in \xi(z, u)}} p(u, z'|z) = 1$$

Let  $p(u|z) := \sum_{z': z' \in \xi(z, u)} p(u, z'|z)$ . In practical application, each prior probability  $p(u, z'|z)$  can be approximated by the ratio of the number of trajectories from  $z$  to  $z'$  over the number of all trajectories from  $z$  to its neighbor states under input  $u$ , which can be obtained by a sufficient number of simulations. Based on this prior probability distribution we can define a local measure of nondeterminism as follows,  $(\forall z \in Z)$

$$H_z := - \sum_{\substack{u, \\ \xi(z, u) \neq \emptyset}} p(u|z) \left[ \sum_{\substack{u, z', \\ z' \in \xi(z, u)}} \frac{p(u, z'|z)}{p(u|z)} \log \frac{p(u, z'|z)}{p(u|z)} \right] \quad (7)$$

We can check that  $H_z$  becomes zero when transitions at  $z$  are deterministic. Figure 7 depicts three different local transition structures with different measures of nondeterminism, where clearly  $0 = H_z^a < H_z^c < H_z^b$ . So

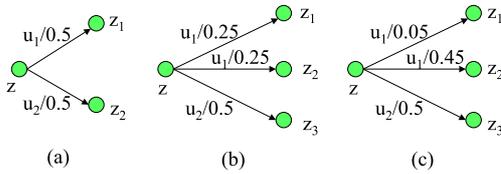


Figure 7: Transition Models with Different Degrees of Nondeterminism

we can conclude that at state  $z$ , transition structure (c) is much closer to the deterministic structure (a) than (b) does.

The global measure of nondeterminism in a discrete-event model can be simply taken as the average over all

local measures, namely

$$H_{ND} := \frac{\sum_{z \in Z} H_z}{|Z|}$$

where  $|Z|$  is the cardinality of  $Z$ . Based on this measure we can show that a model  $DM$  is deterministic if and only if  $H_{ND}^{DM} = 0$ , which is equivalent to that  $(\forall z \in Z) H_z = 0$ . We say a model  $DM_1$  is more nondeterministic than a model  $DM_2$  if  $H_{ND}^{DM_1} > H_{ND}^{DM_2}$ . In other words,  $H_{ND}$  over the degree of nondeterminism of discrete models is nonnegative monotonically increasing.

In the two-tank system let us assume for the sake of example that the prior probability is uniform over *facet-crossing* transitions ( $z_j \in \xi(z_i, u) \Rightarrow \dim(B(z_i, z_j)) = n - 1$ ) and zero over *non-facet-crossing* transitions ( $z_j \in \xi(z_i, u) \Rightarrow \dim(B(z_i, z_j)) < n - 1$ ). The reason for zero-probability assignment on non-facet-crossing transitions is that for each state  $z$ , the ratio of total number of trajectories starting from  $z$  and crossing non-facet boundary over total number of trajectories starting from  $z$  is zero. With such probability assignment, the measure of nondeterminism is  $H_{ND} = 0.414721$ . If we increase the total number of states from 9 to 16 by simply increasing the number of divisions over each tank height from 3 to 4, then  $H_{ND} = 0.461403$ , which indicates that such partition refinement won't reduce the nondeterminism.

This result can be intuitively explained as follows. In the two-tank discrete model, most states usually have eight neighboring states, except boundary states which have fewer neighbors. The measure of nondeterminism of a state with eight neighbors is higher than that of a boundary state. When we increase the number of states by the proposed way, each non-boundary state still has eight neighbors and the transition structure of such a non-boundary state is the same as in the fewer-state model. Since a finer partition means a larger portion of states with eight neighbors, the global measure of nondeterminism increases. On the other hand, when the total number of states become larger and larger, the portion of states with eight neighbors become dominant. So we expect that the global measure of nondeterminism will asymptotically approach to a fixed number. Figure 8 depicts another example where increasing the number of states from (A) to (B) cannot reduce the measure of nondeterminism.

Although using the proposed way to build a model with more states won't reduce the degree of nondeterminism, it may make it possible for us to do supervisory control if the original fewer cell model cannot. This is nothing to do with nondeterminism, but to do with how much details an abstract model can reveal. For example, if we want to maintain the trajectories within the area of  $[3, 9) \times [3, 9)$ , then we cannot model this area with one single state because this will leave no place for us to take control actions. So we have to refine our par-

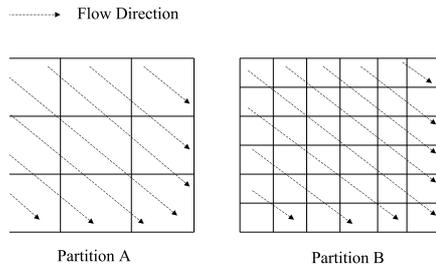


Figure 8: Effect of Partition on Nondeterminism

tion and model this area with more states, which will make it possible for us to take control actions to maintain trajectories within these states, hence within the desirable area. So when two models with similar global measure of nondeterminism, the model with more states is usually better. On the other hand, when we increase the number of states, if the global measure of nondeterminism increases drastically, then it strongly suggests that the current partition method is not good for control purpose. So this entropy measure offers us a way to evaluate the quality of the current partition method. But it cannot suggest us how to improve a partition if the partition is found not good enough. As for how to efficiently refine a partition to reduce nondeterminism, it is still under investigation.

## 5 Conclusions

In this paper we have shown that, with a partition consisting of regular polyhedra, we can generate discrete abstraction of a continuous-time system by using gradient analysis on boundaries of each cell in the partition. It is more computationally efficient than other trajectory-based discrete abstraction methods in the literature as long as a continuous-time system with a regular polyhedron state space is concerned. After a discrete model is obtained, we illustrate how to apply supervisory control on such a model. Then considering that the resulting discrete model is usually nondeterministic, we propose an entropy measure on nondeterminism, which can be used as a quality index of the current partition. There are still several problems left unsolved, e.g. how to expand the proposed method to a system whose domain is not a regular polyhedron, and how to reduce nondeterminism after the entropy index indicates that the current partition is not satisfactory. These problems will be discussed in our future work.

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